

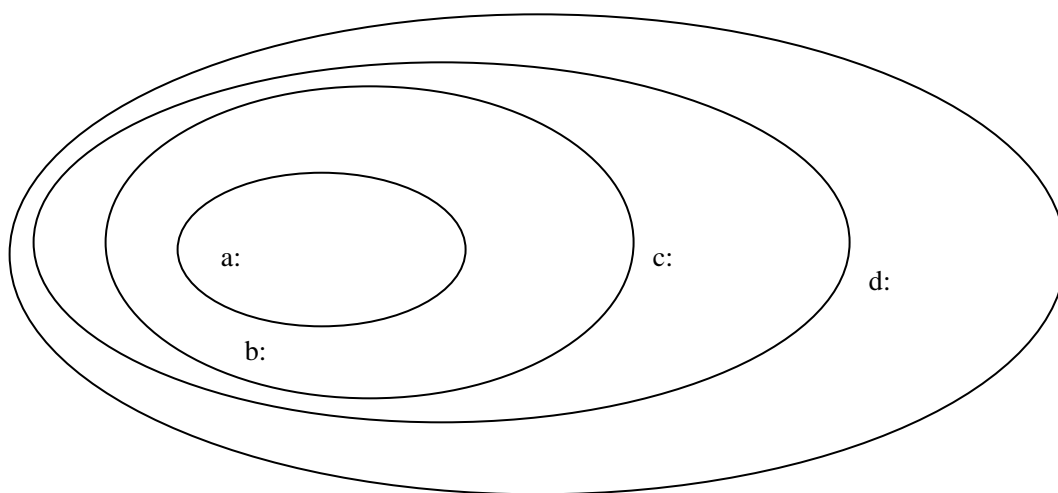
## Theory of Computation Practice Final

1. **True/False questions:** For each part, circle either True or False. (23 points: 1 points each)

- |   |      |       |
|---|------|-------|
| a. A TM can compute anything a desktop PC can, although it might take more time.  | True | False |
| b. A Push-Down Automata can compute things that a TM cannot compute.  | True | False |
| c. Every Turing-decidable language is also Turing-recognizable.   | True | False |
| d. The Halting problem is decidable.  | True | False |
| e. All problems have an algorithm that will solve/decide them.  | True | False |
| f. $4n^2 = O(n)$  | True | False |
| g. $3n^3 = O(n^3)$  | True | False |
| h. $n \log n = o(n^2)$  | True | False |
| i. $2^n = O(n^{20})$  | True | False |
| j. You <u>cannot</u> build a DFA to recognize $\{0^{500}1^{10000} \cup 1^{1000}0^{200}\}$   | True | False |
| k. NP is the class of languages with polynomial time verifiers.   | True | False |
| l. The various sorting algorithms (e.g., bubblesort, heapsort) are in NP  | True | False |
| m. Most theoretical computer scientists believe that $P = NP$ .   | True | False |
| n. All languages are Turing-recognizable  | True | False |
| o. The class of regular languages is closed under union   | True | False |
| p. All languages are decidable  | True | False |
| q. A regular language L may not be context-free.  | True | False |
| r. $A_{DFA}$ is decidable. $A_{DFA} = \{ \langle B, w \rangle \mid B \text{ is a DFA that accepts input string } w \}$ .                                    | True | False |
| s. Deterministic and non-deterministic PDA's have equivalent expressive power.  | True | False |
| t. If a problem A is reducible to problem B, then problem A must be no harder than B.   | True | False |
| u. An algorithm implemented on a single tape Turing machine will always have the same running time (e.g., big-O value) when run on a 2-tape Turing machine. | True | False |
| v. NP is the class of languages decided by some nondeterministic polynomial time Turing Machine.  | True | False |
| w. From a computability perspective, every multi-tape Turing machine has an equivalent single-tape TM.  | True | False |

2. **Short answer questions.** Answer each question in a few sentences. (14 points: 2 each)

- a. The diagram below show a hierarchy of the languages we learned, with respect to computability. Write the proper language next to the labels a-d in the diagram below such that the hierarchy is correct. The languages are: Turing-recognizable, regular, decidable, context-free.



- b. A finite automata will run until its input is completely processed and then it will stop. This is not true for a Turing machine. Explain why.
- c. A language is Turing-recognizable if some Turing machine recognizes it (this is a definition). But what does it mean when we say that a TM recognizes a language? The answer can be quite simple (one sentence) but please be precise.

- d. A language is Turing-decidable if some Turing machine decides it. What does it mean for a Turing machine to decide a language? Again, please be precise, but you can be relatively informal.
- e. We are given a problem and find out that it is undecidable. Could there be an algorithm to solve it in polynomial time? Answer “yes” or “no” and then explain/justify your answer.
- f. I tell you that an algorithm runs in  $O(2^n)$  but yet is in P. How can this be?
- g. Assume that someone finds a polynomial time solution to an NP-complete problem. 1) What does that say about all NP-complete problems? 2) What does that say about the question of whether  $P = NP$ ?



6. (10 points total) Give implementation level descriptions of Turing machines that decide the following languages over the alphabet  $\{a,b,c\}$ . Recall that implementation level is lower level than the pseudo code that we use to describe algorithms (at the implementation level you talk about scanning the tape and tape movements).

$\{w \mid w \text{ contains more than 5 times as many b's as a's}\}$

7. (5 points) Let  $A_{\text{DFA}} = \{ \langle B, w \rangle \mid B \text{ is a DFA that accepts input string } w \}$ . Provide a simple explanation (one could call it a very informal proof) for why this language is decidable.
8. (9 points) Describe a PDA that accepts the following languages. For part b, if you want, rather than describing it from scratch, you can just say how you would modify your answer from part a. Note: part b is a bit tricky. Hint: for part b, you will need to utilize a capability of the PDA not used in part a.
- a.  $L = \{0^m 1^n : n \leq m\}$