

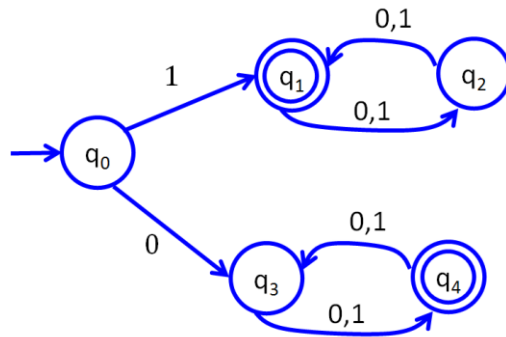
1. Consider the state diagram for the following DFAs. For each, answer the following questions:

(1) What state is reached by the input:  $w=00110$  ?

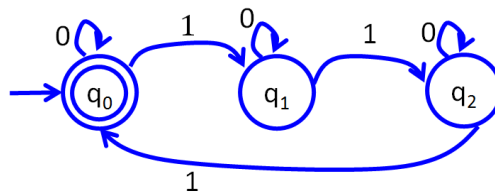
(2) What is the transition function?

(3) What is the language recognized?

**M1:**



**M2:**



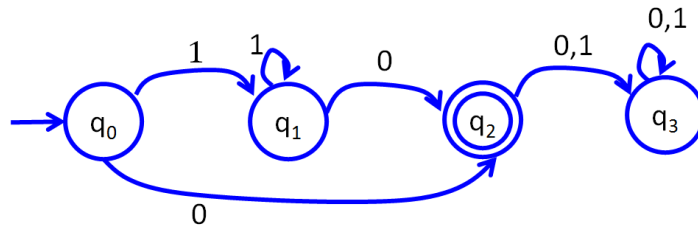
(1)  $q_2$        $q_0(0) \rightarrow q_0(0) \rightarrow q_0(1) \rightarrow q_1(1) \rightarrow q_2(0) \rightarrow q_2$

(2)

	0	1
$q_0$	$q_0$	$q_1$
$q_1$	$q_1$	$q_2$
$q_2$	$q_2$	$q_0$

(3)  $\{w \mid \text{the number of 1's entered is a multiple of 3}\}$   
 $\{0^*(10^*10^*10^*)^*\}$

**M3:**



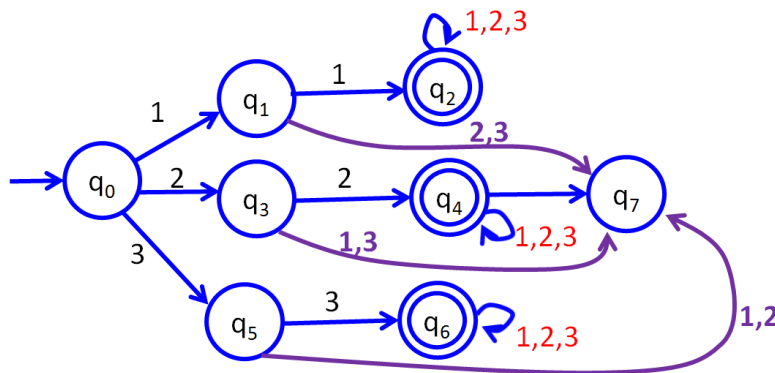
2. Define a machine to recognize the following languages in the alphabet  $\Sigma = \{1,2,3\}$

(5 points)

**L4**={w | the product of input symbols is even} E.g., 111 → 1x1x1=1 is odd-reject, 233 → 2x3x3 = 18 is even-accept

**L5**={w | numbers entered in non-decreasing order} Examples: 112223, 122333

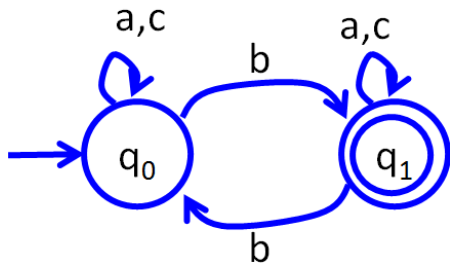
**L6**={w | first two symbols are identical} Examples: 001213, 333212, 3310013



3. Prove the following languages are regular, using the alphabet  $\Sigma = \{a, b, c\}$ :

**L7**={w | w contains an odd number of b's}

Define a DFA to detect the language and/or show a regular expression captures the language.



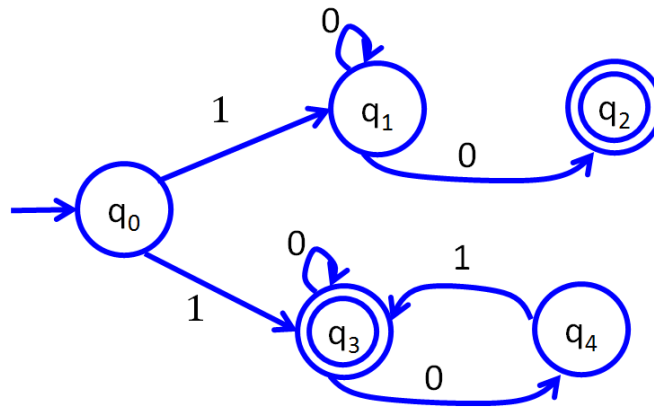
$(aUc)^*b(aUc)^*((aUc)^*b(aUc)^*b(aUc)^*$

$L8 = \{w \mid w \text{ contains the sequence } bcb\}$  (Examples: aabbbcb or cbcba)

$L9 = \{w \mid w \text{ does not have three a's in a row}\}$

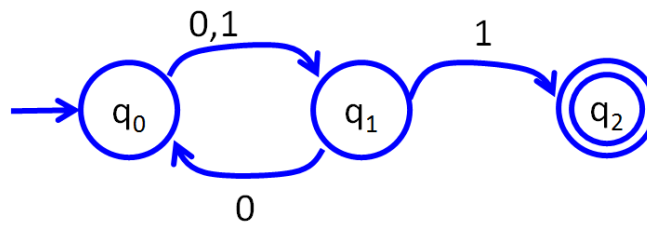
4. Consider the following NFAs. For each, answer:
- (1) what state(s) will be reached by the input: 0011
  - (2) provide a regular expression to describe the recognized language
  - (3) For **N11** and **N12**, convert NFA to DFA

**N10:**

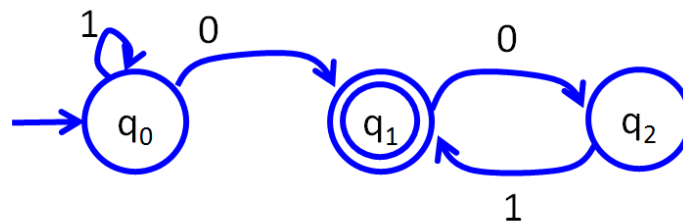


- (1) No state – it will be rejected!
- (2)  $1(0^* \cup 0^*(0^*01)^*$

**N11:**



**N12:**



5. For each regular expression using  $\Sigma = \{a, b\}$ :

(1) Provide three example words.

(2) Convert these regular expressions to a DFA or NFA

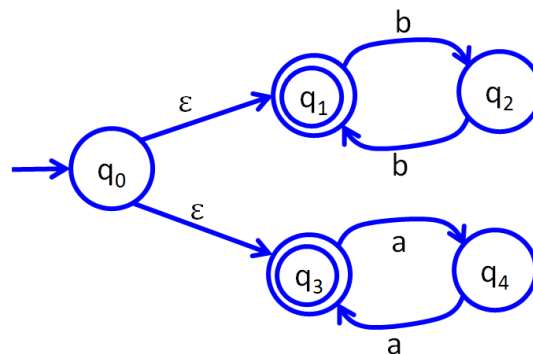
**L13**={**ab\*(ba)\***}

**L14**={(**a ∪ b**)**ba\***}

**L15**={(**bb**)\* ∪ (**aa**)\*}

(1) Examples: **bb, aa, bbbbbb, aaaaaa, ε**

(2)



6. What is the minimum pumping length for each of these languages, showing these languages are regular? We use the alphabet  $\Sigma = \{0,1,2\}$

**L16**={**00(0 ∪ 1)\*12**}

**L17**={**0(22)\*10**}

**p=5**, minimum pumpable string is 02210

**L18**={**111(202)\*210**}

If pumping length is  $p=5$ , how would you break up string  $w$  into  $x$ ,  $y$ , and  $z$  for languages  $L$  below?

$L_{19} = \{ 20(11)^*001 \}$ ,  $w=201111001$

$L_{20} = \{ (121)^*001 \}$   $w=121001$

7. Consider the language  $L_{21} = \{01(101)^*11\}$ , what is the error in each of the following “pumping lemma” arguments?

**Argument 1:** Let us take  $w=0111$ ,  $w \in L_{21}$ . We cannot divide  $w=xyz$  such that  $y^i z \in L_{21}$ ,  $i \geq 0$ . For example, if  $x=0$ ,  $y=11$ , and  $z=1$ ,  $xy^2z = 011111 \notin L_{21}$ . Therefore,  $L_{21}$  is not regular.

The pumping length is  $p=7$ . Using any strings in  $L_{21}$  with length less than pumping length is not necessarily pumpable, and the inability to pump a too-short string does not prove anything. You can only test pumping on strings with at least as many characters as the pumping length.

**Argument 2:** Let us take  $w=0110110111$ ,  $w \in L_{21}$ . If we divide  $w=xyz$  as follows:  $x=0110110$ ,  $y=11$ ,  $z=1$ , we cannot repeat  $y$  such that  $xy^i z \in L_{21}$ ,  $i \geq 0$ . For example, if  $xy^2z = 011011011111 \notin L_{21}$ . Therefore,  $L_{21}$  is not regular.

8. Prove these languages are not regular.

$$L24 = \{0^n 1^{2n} 0^{3n} \mid n > 0\}$$

Proof by contradiction with pumping lemma:

Assume L24 is pumpable. Now consider  $w = 0^p 1^{2p} 0^{3p}$ , which is element of L24 with  $|w| > p$ . Thus,  $w$  must be pumpable.

$$w = xyz \quad x = 0^j \quad y = 0^k \quad z = 0^{p-(j+k)} 1^{2p} 0^{3p} \quad j+k \leq p$$

$$\text{Try pumping } w: \quad xy^2z \rightarrow 0^j 0^{2k} 0^{p-(j+k)} 1^{2p} 0^{3p}$$

$xy^2z$  begins with  $j+2k+p-(j+k)$  0's ...  $j+2k+p-(j+k) = p+k$  0's

$xy^2z$  begins with  $p+k$  0's followed by  $2p$  1's.

$2p \neq p+k$ , so  $xy^2z \notin L24$ , which means L24 is not regular!

$$L25 = \{1^{n^3} \mid n > 0\}$$

9. For each of the following grammars, list three strings produced by the grammar

**G26:**

$$S \rightarrow AB \mid BA$$

$$A \rightarrow xAy \mid \epsilon$$

$$B \rightarrow BzB \mid y$$

$$\text{Examples: } AB \rightarrow \epsilon y \rightarrow y$$

$$BA \rightarrow yxAy \rightarrow yxxAyy \rightarrow yxx\epsilon yy \rightarrow yxxyy$$

$$AB \rightarrow \epsilon BzB \rightarrow \epsilon BzBzy \rightarrow \epsilon BzBzyzy \rightarrow \epsilon yzyzyzy$$

**G27:**

$$S \rightarrow A \mid AA$$

$$A \rightarrow 00 \mid 11$$

**G28:**

$$A \rightarrow 11A00 \mid \epsilon$$

10. Provide the languages described by two of the grammars:

**G27 (from above)**

**G28 (from above)**

10. Provide a grammar to produce the following languages

**L32 =  $\{0^n(11)^n \mid n \geq 0\}$**

**L33 =  $\{01^*00^*\}$**

**L34 =  $\{w \mid w = w^{\text{Reverse}}\}$     Examples: 00100, 10101, 1111**

**S  $\rightarrow$  0S0 | 1S1 | 0 | 1 |  $\epsilon$**

11. Convert the following grammars to Chomsky Normal Form

**G29:**

**S  $\rightarrow$  xAy | BA**

**A  $\rightarrow$  z | AzA**



**$B \rightarrow yB \mid \epsilon$**

Remove the  $\epsilon$

**$S \rightarrow xAy \mid BA \mid A$**

**$A \rightarrow z \mid AzA$**

**$B \rightarrow yB \mid y$**

Remove the  $S \rightarrow A$  unit rule

**$S \rightarrow xAy \mid BA \mid z \mid AzA$**

**$A \rightarrow z \mid AzA$**

**$B \rightarrow yB \mid y$**

Replace mixed terminal—variable rules with variables-only rules

**$S \rightarrow U_xAU_y \mid BA \mid z \mid AU_zA$**

**$U_x \rightarrow x$**

**$U_y \rightarrow y$**

**$U_z \rightarrow z$**

**$A \rightarrow z \mid AU_zA$**

**$B \rightarrow U_yB \mid y$**

Replace 3+ variable rules with 2-variables rules

**$S \rightarrow U_xC \mid BA \mid z \mid AD$**

**$C \rightarrow AU_y$**

**$D \rightarrow U_zA$**

**$U_x \rightarrow x$**

**$U_y \rightarrow y$**

**$U_z \rightarrow z$**

**$A \rightarrow z \mid AD$**

**$B \rightarrow U_yB \mid y$**

**G30:**

**$S \rightarrow BAB \mid ABA$**

**$A \rightarrow y \mid z$**

$B \rightarrow x \mid AA \mid \varepsilon$

**G31:**

$S \rightarrow ByBy$

$B \rightarrow xBx \mid \varepsilon$